

# Metalogica

Ad constructione de apparatus de Metalogica  
 que es generalizatione de logica. In isto  
 ubi symbolas  $a, b, \dots$  habe duo valore: classe et  
 propositione (in Metalogica hyperclasse et hyper  
 propositione) licet ut loco de duo operatione  
 fundamentale: de additione et multiplicatione  
 caret tres operationes cum notationes:

$$abc, \overline{abc}, \underline{abc}$$

loco de symbolos: veri  $\vee$  et falso  $\wedge$

tres symbolos:  $\vee, \overline{\vee}, \underline{\vee}$

loco duo termino  $a, a$  genere et specie, propositione  
 $a \vee a$  et  $a \wedge a$  - tres terminos specie  $a, a, a$ ,  
 hypergenere  $a''$ , hyp. dato  $a$ , educto  $a'$ , hypereducto  $a''$

Es necesse etiam de postulato.

$$a \vee \overline{a} = \vee$$

$$\overline{a \vee \overline{a}} = \overline{\vee}$$

$$\underline{a \vee \overline{a}} = \underline{\vee}$$

## commutativo

$$abc = bca = cab$$

$$\overline{abc} = \overline{bca} = \overline{cab}$$

$$\underline{abc} = \underline{bca} = \underline{cab}$$

## associativo

$$ab(cde) = (abc)de = b(cde) = (d(eab)) = de(ab) = e(abd)$$

In omnes hoc leges es symbolo  $\equiv$  de identitate que

est disinpretatio a valitate circulari

$$a = b = \dots$$

que reduce ad identitate cum  $b \equiv a$

in omnes se repetit te circulari:

$$|a''b'c| > |b''c'a| > |c''a'b|$$

Alios nos ulos:

de simplicitate:

$$|abc'' a'a| \& |a'' abc' a| \& |a'' a' abc|$$

de compositione

$$\left| \begin{array}{l} a' b' c \\ a'' d' e \\ a'' f' n \end{array} \right| \supset \left| a'' b' d' f' e' h \right| \left| \begin{array}{l} a'' b' c \\ d'' e' f' \\ f'' b' n \end{array} \right| \supset \left| a' d' f'' e' c' h \right|$$

$$\left| \begin{array}{l} a'' b' c \\ a'' e' f' \\ f'' h' c \end{array} \right| \supset \left| a' d' f'' b' e' h', c \right|$$

de leutologia

$$aaa \equiv a, \widehat{aaa} \equiv a, \underline{aaa} \equiv a$$

Nos da duo exemplo de deductione

Postulato de hypersyllogismo:

$$\left| \begin{array}{l} a'' b' c \\ d'' n' i \\ g'' d' a \end{array} \right| \supset |g'' i' c|$$

Nos da uno exemplo de deductione:

Teorema que responde ad teorema de multiplicatione:

$$|a'' b' c| \supset |a' d' g'' b' d' e' f' c' d' e' g'|$$

secundu postulato de hypersyllogismo

$$\left| \begin{array}{l} a'' b' c \\ a'' b' c \\ a' d' g'' a' a \end{array} \right| \supset |a' d' g'' b' c|$$

et nos tunc de compositione

$$\left| \begin{array}{l} a' d' g'' b' c \\ a' d' e' f' i' i \\ a' d' e' f' g' g \end{array} \right| \supset |a' d' g'' b' d' e' f' c' d' e' g'|$$

Exemplo secundo - deductione formula de

absumptio e

$$\underline{abc} . aa \equiv a$$

Per postulato de simplicatione

$$|a'' a' \underline{abc . aa}|$$

de compositione

$$\left| \begin{array}{l} a'' a' a \\ a'' a' a \\ abc'' a' a \end{array} \right| \supset |abc . aa'' a' a|$$

$$\left| \begin{array}{l} a'' a' a \\ a'' a' a \\ a'' a' a \end{array} \right| \supset |a'' abc . aa' a|$$

unde sequitur equalitate

$$\underbrace{abc} \cdot aa = a = a$$

et identitate

$$\underbrace{abc} \cdot aa \equiv a$$

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## Meta algebra

Ad constructionem de Meta algebra nos debet  
capere loco duo operationes fundamentales:  
additione et multiplicatione tres operationes  
super tres objecto cum notationes:

$$abc, \widehat{abc}, \underline{abc}$$

Leges fundamentales:

### commutativo

$$\text{I} \quad abc \equiv bca \equiv cab$$

$$\widehat{abc} \equiv \widehat{bca} \equiv \widehat{cab}$$

$$\underline{abc} \equiv \underline{bca} \equiv \underline{cab}$$

### associativo

$$(abc)de \equiv (bcd)ea \equiv (cde)ab$$

$$\widehat{(abc)de} \equiv \widehat{(bcd)ea} \equiv \widehat{(cde)ab}$$

$$\underline{(abc)de} \equiv \underline{(bcd)ea} \equiv \underline{(cde)ab}$$

### distributivo de primo specie

$$\widehat{a}gh \cdot \widehat{b}gh \cdot \widehat{c}gh = \widehat{abc} \cdot gh$$

$$\underline{a}gh \cdot \underline{b}gh \cdot \underline{c}gh = \underline{abc} \cdot gh$$

$$\widehat{a}gh \cdot \widehat{b}gh \cdot \widehat{c}gh = \widehat{abc} \cdot \widehat{gh}$$

Differentia essentialis in comparatione cum

Meta algebra constat in hoc: secundum et tertium

operationes et distributivum respectu ad primum

sed respectu ad secundum solum tertium et distributivum

Hoc legem respondet ad formulam

$$(a + b)y = ay + by$$

distributivo de secondo specie

$$\text{IV) } \begin{aligned} abc \cdot def \cdot ghi &\equiv acd \cdot beh \cdot cfi \\ \overline{abc \cdot def \cdot ghi} &\equiv \overline{acd \cdot beh \cdot cfi} \\ \underline{abc \cdot def \cdot ghi} &\equiv \underline{acd \cdot beh \cdot cfi} \end{aligned}$$

hoc es analogo de formula  
 $ab \cdot cd = ca \cdot bd$   
 $(a+b) + (c+d) = (a+c) + (b+d)$

que in alio potest demonstrare per I, II, III  
 sed hic remaneat de postulata

$$\text{V) } \begin{aligned} \overline{abc \cdot def} \cdot ghi &= a \cdot b \cdot df \cdot e \cdot gh \\ \overline{abc \cdot de} \cdot f \cdot ghi &= a \cdot b \cdot f \cdot e \cdot gh \\ \overline{abc \cdot de} \cdot f \cdot ghi &= a \cdot b \cdot df \cdot e \cdot gh \end{aligned}$$

VI) associativo

$$\overline{abc} \cdot de \equiv$$

Actone de separate formulae de Mathematica  
 posse de vide in theoria de hyperfractiones  
 Fractione posse de definir per equatione

$$x \alpha \equiv a$$

ubi a numeratore,  $\alpha$  denominatore

In hoc equatione ~~est~~ substitue

$$x \alpha \alpha' \equiv a$$

Hyperfractione habe uno numeratore et duo  
 denominatore

Nos i ostende quomodo posse applica leges  
 fundamentales ad deductione ~~ita~~ regulas  
 de additione + multiplicatione de hyperfractione

$$\begin{array}{l|l} x \alpha \alpha' \equiv a \\ y \alpha \alpha' \equiv b \\ z \alpha \alpha' \equiv c \end{array} \left| \begin{array}{l} \overline{x \alpha \alpha' \cdot y \alpha \alpha' \cdot z \alpha \alpha'} = \overline{abc} \\ \text{et per (III)} \\ \underline{xyz \cdot \alpha \alpha'} = \underline{abc} \end{array} \right.$$

Et dem modo

$$\frac{a}{\alpha|\alpha'} \cdot \frac{b}{\alpha|\alpha'} \cdot \frac{c}{\alpha|\alpha'} \equiv \frac{abc}{\alpha|\alpha'}$$

In casu de denominatores diversos  
necesse es inveniuntione de maximo commune  
multiplo.

Invenito

$$\begin{array}{c|c|c} A_1 A_2 & B_1 B_2 & C_1 C_2 \\ \hline A'_1 A'_2 & B'_1 B'_2 & C'_1 C'_2 \end{array}$$

tales, ut

$$\alpha A_1 A_2 \equiv \beta B_1 B_2 \equiv \gamma C_1 C_2 = M$$

$$\alpha' A'_1 A'_2 \equiv \beta' B'_1 B'_2 \equiv \gamma' C'_1 C'_2 = M'$$

Multiplicato  $x \alpha \alpha' \equiv a$  per  $A_1 A_2$

$$(x \alpha \alpha') A_1 A_2 \equiv a A_1 A_2$$

et postea per  $A'_1 A'_2$

$$[(x \alpha \alpha') A_1 A_2] A'_1 A'_2 \equiv (a A_1 A_2) A'_1 A'_2$$

et per (V)

$$x (\alpha A_1 A_2) (\alpha' A'_1 A'_2) \equiv (a A_1 A_2) A'_1 A'_2$$

$$x \equiv \frac{(a A_1 A_2) \cdot A'_1 A'_2}{M M'}$$

$$y \equiv \frac{(b B_1 B_2) \cdot B'_1 B'_2}{M M'}$$

$$z \equiv \frac{(c C_1 C_2) \cdot C'_1 C'_2}{M M'}$$

Nos vide ut additione de hyperfractione cum  
denominatores diversos redi ad additio e de  
hyperfractione cum denominatores identicos

Multiplicatione de hyperfractiones facto per  
hoc modo

$$\left. \begin{array}{l} x \alpha \alpha' \equiv a \\ y \beta \beta' \equiv b \\ z \gamma \gamma' \equiv c \end{array} \right\} x \alpha \alpha' \cdot y \beta \beta' \cdot z \gamma \gamma' \equiv abc$$

Et per (IV)  $x y z \cdot \alpha \beta \gamma \cdot \alpha' \beta' \gamma' = abc$

$$\frac{a}{\alpha|\alpha'} \cdot \frac{b}{\beta|\beta'} \cdot \frac{c}{\gamma|\gamma'} \equiv \frac{abc}{\alpha\beta\gamma|\alpha'\beta'\gamma'}$$

Operatione de subtractione et de numeros  
negativos definito per equatione

$$\widehat{x\alpha} \equiv a$$

de que analogo

$$\widehat{x\alpha\alpha'} \equiv a$$

aut

$$\widehat{x\alpha\alpha'} \equiv a$$

unde evolve teoria de hypernumeros negativos.

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